# An Insertion Algorithm for Multiplying Schubert Polynomials by Schur Polynomials

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# Outline

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#### Schubert polynomials

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Grassmannian permutations Lattice chains in *k*-Bruhat order Lattice Permutation Tableaux

# A Littlewood–Richardson Rule

for Grassmannian permutations

Schubert structure constants  $c_{u,v}^w$ , triply indexed by permutations, defined by

$$\mathfrak{S}_{u}\mathfrak{S}_{v}=\sum_{w}c_{u,v}^{w}\mathfrak{S}_{w},$$

enumerate flags in a suitable triple intersection of Schubert varieties, and so  $c_{u,v}^w \in \mathbb{N}$ .

**Fundamental open problem**: Give a simple positive combinatorial formula for  $c_{u,v}^{w}$ .

A permutation v is k-grassmannian if  $v_i < v_{i+1}$  for all  $i \neq k$ ; e.g. 1367245 is 4-grassmannian.

A *k*-partition is a weakly increasing sequence of length k; e.g. (0, 1, 3, 3) is a 4-partition.

The map  $v \mapsto v(v, k)$  given by  $v_i = v_i - i$  for  $i \le k$  is a bijection between *k*-grassmannian permutations and *k*-partitions; e.g. 1367245 = v((0, 1, 3, 3), 4).

My conjectured formula, proved with Nantel Bergeron, states for all u, v, k, w,

$$c_{u,\nu(\nu,k)}^{w} = \# \begin{cases} \text{saturated chains in } k\text{-Bruhat order} \\ \text{from } u \text{ to } w \text{ with lattice weight } \nu \end{cases}$$

Grassmannian permutations Lattice chains in *k*-Bruhat order Lattice Permutation Tableaux

# k-Bruhat Order

# Lattice chains

Bruhat order has cover relations  $u < ut_{r,s}$  if  $\ell(ut_{r,s}) = \ell(u) + 1$ , where  $\ell(u) = \#\{i < j \mid u_i > u_j\}$ .

#### Definition (Bergeron–Sottile 1998)

 $u \leq_k w$  is the transitive closure of  $u \prec_k ut_{a,b}$  whenever  $a \leq k < b$  with  $\ell(ut_{a,b}) = \ell(u) + 1$ 

Bergeron and Sottile decorate saturated chains by labeling the cover  $u \prec_k ut_{a,b}$  with  $u_b$ .

 $136\underline{2}\underline{5}847 \xrightarrow{5} 13\underline{6}5|284\underline{7} \xrightarrow{7} 13\underline{7}5|2\underline{8}46 \xrightarrow{8} \underline{1}385|\underline{2}746 \xrightarrow{2} 2\underline{3}85|17\underline{4}6 \xrightarrow{4} 248\underline{5}|1\underline{7}36 \xrightarrow{7} 2\underline{4}87|1\underline{5}36 \xrightarrow{5} 25871436$ For  $d_1, \dots, d_m$  the decorations of C, set  $e_i = \begin{cases} k & \text{if } i = 1 \\ e_{i-1} & \text{if } d_{i-1} < d_i \\ e_{i-1} - 1 & \text{if } d_{i-1} > d_i \end{cases}$ The weight is wt $(C)_i = \{j \mid e_j = i\} = (0, 1, 3, 3)$   $\begin{pmatrix} 5 & 7 & 8 & 2 & 4 & 7 & 5 \\ 4 & 4 & 4 & 3 & 3 & 2 \end{pmatrix} \xrightarrow{\text{sort top}} \begin{pmatrix} 2 & 4 & 5 & 5 & 7 & 7 & 8 \\ 3 & 3 & 4 & 2 & 4 & 3 & 4 \end{pmatrix}$ 

#### Definition (A. 2021)

A chain C is lattice if for all d, for all  $i \le k$ , we have  $\#\{j \mid e_j = i - 1, d_j \ge d\} \le \#\{j \mid e_j = i, d_j \ge d\}$ .

#### Theorem (A.–Bergeron 2023)

 $c_{u,v(\nu,k)}^{w} = \#\{\text{saturated chains in } k\text{-Bruhat order from } u \text{ to } w \text{ with } \text{lattice weight } \nu\}$ 

Grassmannian permutations Lattice chains in *k*-Bruhat order Lattice Permutation Tableaux

# Lattice Permutation Tableaux



#### Definition (A. 2021)

For C from u to w, set  $T_0 = \mathbb{D}(u)$ , and define  $T_{i+1}$  by

- move cells in row  $b_i$  right of column  $u_{a_i}^{(i)}$  down to row  $a_i$
- place a new skew cell with entry  $e_i$  in row  $a_i$ , column  $u_{a_i}^{(i)}$

between a<sub>i</sub>, b<sub>i</sub> and u<sup>(i)</sup><sub>a<sub>i</sub></sub>, u<sup>(i)</sup><sub>b<sub>i</sub></sub>, move cells left to avoid dots



#### Definition (A. 2021)

LPT<sub>k</sub>( $w/u, \nu$ ) is the set of permutation tableaux from u to w in k-Bruhat with lattice weight  $\nu$ .

Kohnert's rule Pipe dreams Main Theorem

# Kohnert's rule

#### Definition (Kohnert 1991)

A Kohnert move on a row selects the rightmost cell c, moves c down to the first open position below.

 $KD(D) = \{T \mid T \text{ obtained by Kohnert moves on } D\}$ 



#### Theorem (Kohnert 1991)

The Schur polynomial indexed by  $\nu$  is

$$s_{\nu}(x_1,\ldots,x_k) = \sum_{T \in \mathrm{KD}(\mathbb{D}(\nu(\nu,k)))} x_1^{\mathrm{wt}(T)_1} \cdots x_n^{\mathrm{wt}(T)_n}$$

Placing an r into cells in row r and floating up to row k gives a weight-preserving bijection

 $\mathrm{KD}(\mathbb{D}(v(\nu,k))) \xrightarrow{\varphi} \mathrm{SSRT}_k(\nu)$ 

with semi-standard reverse tableaux.

#### Theorem (Winkel 1999, Winkel 2002, A. 2022, Armon–A.–Bowling–Ehrhard 2023)

The Schubert polynomial indexed by w is  $\mathfrak{S}_{w} = \sum_{T \in \mathrm{KD}(\mathbb{D}(w))} x_{1}^{\mathrm{wt}(T)_{1}} \cdots x_{n}^{\mathrm{wt}(T)_{n}}$ 

AABE 2023: Kohnert moves generate the character of any southwest flagged Schur module.

Pipe dreams

# Kohnert diagrams

Theorem (Billey–Jockush–Stanley 1992, Bergeron–Billey 1993)

The Schubert polynomial  $\mathfrak{S}_w = \sum_{P \in \mathbb{RPD}(w)} x_1^{\operatorname{wt}(T)_1} \cdots x_n^{\operatorname{wt}(T)_n}$  is generated by reduced pipe dreams.

Pipe dreams

(1) Remove all  $\checkmark$ s (2) Shift +s in row r right r-1 columns (3) Change + to  $\bigcirc$  and rectify



To rectify at (c, r), pair cells  $\{(c+1, s) | s > r\}$  to the nearest unpaired cell  $\{(c, s) | s > r\}$  weakly above, if it exists, and then move all unpaired cells in column c + 1 left to column c.

Theorem (A. 2022)

This is a weight-preserving bijection between reduced pipe dreams and Kohnert diagrams.

# A bijective proof

Recall RSK gives a bijection  $SSRT_k(\mu) \times SSRT_k(\nu) \xrightarrow{\sim} \bigsqcup_{\mu \subset \lambda} SSRT_k(\lambda) \times LRT_k(\lambda/\mu, \nu)$ 

#### Theorem (A.–Bergeron 2023)

The k-insertion algorithm on Kohnert diagrams gives a weight-preserving bijection

$$\mathrm{KD}(u) \times \mathrm{SSRT}_k(\nu) \xrightarrow{\sim} \bigsqcup_{u \leq kW} \mathrm{KD}(w) \times \mathrm{LPT}_k(w/u, \nu).$$



Corollary (A.–Bergeron 2023)

$$\mathfrak{S}_{u}s_{\nu}(x_{1},\ldots,x_{k})=\sum_{u\prec_{k}w}\#\mathrm{LPT}_{k}(w/u,\nu)\mathfrak{S}_{w}$$



Labeling Kohnert diagrams Landing and bumping Row-Bumping Lemma

### Insertion on Kohnert diagrams



$$\operatorname{KD}(\mathbb{D}(v(\nu,k))) \xrightarrow{\varphi} \operatorname{SSRT}_k(\nu)$$



For  $T \in KD(v(\mu, k))$  and x in row  $r \leq k$ ,

 $\varphi(\operatorname{rectify}(T \sqcup x)) = \varphi(T) \xleftarrow{\operatorname{RSK}} r$ 





Rectification is well-defined on any T and well-behaved on  $T \in KD(D)$  when D is southwest.

- Question 1: At which cells  $x \notin T$  can we land?
- Question 2: Which cells  $y \in T$  can  $x \notin T$  bump?

Answer: We may *k*-land at and *k*-bump a cell if and only if it has semi-proper label at most *k*.

Littlewood–Richardson Rule Schubert polynomials Insertion Algorithm Row-Bumping Lemma

# Semi-proper labelings

A.–Searles (JCT-A 2018) define the proper labeling of cells of a diagram  $T \in KD(\mathbb{D}(\alpha))$ . A. (Trans. AMS 2022) generalized this to (semi-)proper labelings of cells of  $T \in KD(\mathbb{D}(u))$ .

Given a diagram D, the super-standard labeling puts r into cells in row r.



Using label rectification, we can (attempt to) properly label T w.r.t. D explicitly.

# Theorem (A. 2022)For T and D, the following are equivalent:• $T \in KD(D)$ • T can be properly labeled w.r.t. D• T has a semi-proper labeling w.r.t. D

RSK on Kohnert diagrams Landing and bumping Row-Bumping Lemma

# Landing columns

For  $u \prec ut_{r,s}$ , construct  $\mathbb{D}_0(u, (u_s, r)) \in \mathrm{KD}(u)$ by moving cells in row *s* down to row *r*.

Set 
$$\mathbb{D}(u, (u_s, r)) = \mathbb{D}_0(u, (u_s, r)) \sqcup (u_s, r)$$
.

Theorem (A.-Bergeron 2023)

We have a weight-preserving bijection  $\bigcup_{u \prec_k u t_{r,s}} \mathrm{KD}(\mathbb{D}(u, (u_s, r))) \xrightarrow{\sim} \bigsqcup_{u \prec_k u t_{r,s}} \mathrm{KD}(u t_{r,s})$ 



#### Definition (A.-Bergeron 2023)

For  $T \in KD(u)$  and  $x \notin T$  in column  $u_s$  with s > k, say x is a k-landing spot for T if

 $T \cup x \in \mathrm{KD}(\mathbb{D}(u, (u_s, r)))$  and  $\mathcal{L}(x) = r$  for some semi-proper  $\mathcal{L}$ 



Littlewood–Richardson Rule Schubert polynomials Insertion Algorithm Row-Bumping Lemma

# **Bumpable cells**

#### Definition (A.-Bergeron 2023)

For  $T \in KD(u)$  and  $x \notin T$ , say  $y \in T$  below x is k-bumpable if there exists semi-proper  $\mathcal{L}$  s.t.  $\mathcal{L}$  w.r.t.  $\mathbb{D}(u)$  and  $\mathcal{L}(y) \leq k$  or  $\mathcal{L}$  w.r.t.  $\mathbb{D}_0(u, (u_s, r))$  and  $\mathcal{L}(y) = r$ and the labeling on  $T \sqcup x - y$  obtained by swapping x and y is semi-proper.



#### Definition (A.-Bergeron 2023)

For  $T \in KD(\mathbb{D}(u))$  and  $x \notin T$ , the *k*-insertion of *x* into *T*, denoted by  $T \xleftarrow{k} x$ , is

- IF x is a k-landing spot for T, THEN RETURN  $T \cup x$ ;
- **2** ELSE IF there exists a lowest *k*-bumpable cell *y*, THEN RETURN  $(T \sqcup x y) \leftarrow y$ ;
- **3** ELSE RETURN  $T \xleftarrow{k}{\leftarrow} z$  for rightmost  $z \notin T$  left of x.

Littlewood–Richardson Rule Schubert polynomials Insertion Algorithm Row-Bumping Lemma

# **Row-Bumping Lemma**

#### Lemma (A.-Bergeron 2023)

For  $S \in KD(u)$  and  $i, j \le k$ , let  $T = (S \xleftarrow{k} i) \in \mathbb{D}(ut_{r,s}) = \mathbb{D}(v)$  and  $U = (T \xleftarrow{k} j) \in \mathbb{D}(vt_{q,t})$ . Then if i < j, then  $u_s > v_t$  and if  $i \ge j$ , then  $u_s < v_t$ .



Insert  $R \in SSRT_k(\nu)$  into  $T \in KD(u)$  by *k*-inserting the Yamanouchi work for *R*.

$$\begin{pmatrix} 4 & 4 & 3 & 3 & 3 & 2 \\ 2 & 2 & 1 & 4 & 3 & 2 & 4 \end{pmatrix} \xrightarrow{\text{RSK}} \begin{array}{c} 4 & 4 & 4 & 2 \\ \hline 3 & 2 & 1 \\ 2 & 2 & 2 & 2 \\ \hline 2 & 2 & 2 & 2 \\ \hline \end{array} \times \begin{array}{c} 4 & 4 & 4 & 4 \\ \hline 3 & 3 & 3 & 3 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

#### Lemma (A.–Bergeron 2023)

For  $T \in KD(u)$  and  $R \in SSRT_k(\nu)$ , the recording tableau of  $T \stackrel{k}{\leftarrow} R \in LPT_k(w/u, \nu)$  for  $u \prec_k w$ .

RSK on Kohnert diagrams Landing and bumping Row-Bumping Lemma

# **Reversing insertion**



#### Theorem (A.–Bergeron 2023)

The k-insertion algorithm on Kohnert diagrams gives a weight-preserving bijection

$$\mathrm{KD}(u) \times \mathrm{SSRT}_k(\nu) \xrightarrow{\sim} \bigsqcup_{u \prec_k w} \mathrm{KD}(w) \times \mathrm{LPT}_k(w/u, \nu).$$



# References on arXiv

### Nantel Bergeron and Frank Sottile.

Schubert polynomials, the Bruhat order, and the geometry of flag manifolds. *Duke Math. J.*, 95(2):373–423, 1998.



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Sami Assaf and Nantel Bergeron. An insertion algorithm for multiplying Schubert polynomials by Schur polynomials. arXiv:*coming soon*!

# Thank You